

Amplitudes, Ansätze and Algebraic Geometry

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based on [\[2203.04269\]](#)



Introduction

Multi-scale, Multi-loop Amplitudes

- Loop amplitudes are computed via the master integral decomposition.

$$\mathcal{A}^{(l)} = \sum_k \underbrace{\mathcal{C}_k}_{\text{rational functions}} \times \underbrace{\mathcal{I}_k}_{\text{master integrals}}.$$

- Let's focus on difficulties when computing the **rational functions**.
(Talks on integrals by Anatonela, Claude, Christoph, David, Stefan)
- Chosen framework: use **finite-field evaluations** to determine \mathcal{C}_k .

Numerical Data: $\left\{ \mathcal{C}_k(p_1^{(1)}, \dots, p_n^{(1)}), \dots, \mathcal{C}_k(p_1^{(N)}, \dots, p_n^{(N)}) \right\}$

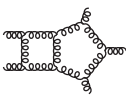
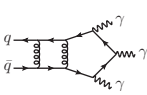
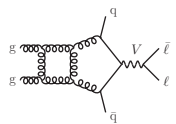
$$\text{Ansatz: } \mathcal{C}_k(p_1, \dots, p_n) = \sum_{i=1}^N c_{ik} a_{ik}(p_1, \dots, p_n)$$

reconstruct \mathcal{C}_k

[von Manteuffel, Schabinger '14; Peraro '16],
FiniteFlow [Peraro '19], Firefly [Klappert, Lange, Klein '19, '20]

Analytic Reconstruction as it Stands

- Reconstruction time **dominated by sampling**: $T_{\text{sample}} \sim O(\text{minute})$.
- Evaluation count** for (selected) recent two-loop five-point amplitudes:

Process			
# Samples	$\sim 10^5$	$\sim 10^5$	$\sim 10^6$.

* After simplification via [Badger et al '20]

Need to better **understand rational functions**, build simpler Ansätze.

The Approach of [De Laurentis, BP '22]

- Work in spinor space* to manifest gauge-theory simplifications.

$$\mathcal{C}_k(p_1, \dots, p_n) \rightarrow \mathcal{C}_k(\lambda, \tilde{\lambda}).$$

*Algorithmic toolkit provided.

- Numerically study \mathcal{C}_k to understand partial-fractions structure.

$$\frac{\mathcal{N}}{\mathcal{D}_1 \mathcal{D}_2 \mathcal{D}_{\text{rest}}} = \frac{\cancel{\Delta}}{\cancel{\mathcal{D}_1} \mathcal{D}_2 \mathcal{D}_{\text{rest}}} + \frac{\Delta_1}{\mathcal{D}_2 \mathcal{D}_{\text{rest}}} + \frac{\Delta_2}{\mathcal{D}_1 \mathcal{D}_{\text{rest}}}$$

See also [De Laurentis, Maître '19].

- Construct Ansatz \mathbf{a}_l from study. Constrain c_{kl} by finite field sampling.

$$\mathcal{C}_k(\lambda, \tilde{\lambda}) = \sum_{l=1}^{N_{\text{new}}} c_{kl} \mathbf{a}_l(\lambda, \tilde{\lambda}), \quad c_{kl} \in \mathbb{Q}, \quad N_{\text{new}} \ll N.$$

The Method, By Example

A First Attempt at Numerical Partial Fractions

- Consider tree-level six-point one-quark line amplitude $\mathcal{A}_{q^+g^+g^+\bar{q}^-g^-g^-}$

$$\mathcal{A} = \frac{\mathcal{N}^*}{\langle 12 \rangle \langle 23 \rangle \langle 34 \rangle [45][56][61] s_{345}}.$$

* \mathcal{N} is a degree 6 **polynomial** in spinor brackets.

- Can we rewrite without **both** $\langle 12 \rangle$ and $\langle 23 \rangle$ poles?

$$\mathcal{A} = \frac{\Delta_{12}}{\langle 23 \rangle \langle 34 \rangle [45][56][61] s_{345}} + \frac{\Delta_{23}}{\langle 12 \rangle \langle 34 \rangle [45][56][61] s_{345}}?$$

- [De Laurentis, Maître '19]: Probe \mathcal{A} on points where $\langle 12 \rangle$, $\langle 23 \rangle$ are small.

$$\lambda_2^\alpha \sim \epsilon \quad \Rightarrow \quad \mathcal{A} \sim \epsilon^{-2}$$

$$\langle 12 \rangle \sim \langle 23 \rangle \sim \langle 13 \rangle \sim \epsilon \quad \Rightarrow \quad \mathcal{A} \sim \epsilon^{-1}.$$

Thinking in Terms of Polynomials

- Let's ask an **equivalent question**:

$$\mathcal{N} = \Delta_{12}\langle 12 \rangle + \Delta_{23}\langle 23 \rangle?$$

- Mathematically, we can ask if \mathcal{N} belongs to an **"ideal"**:

$$\mathcal{N} \in \langle \langle 12 \rangle, \langle 23 \rangle \rangle?$$

- Ideal is infinite set of polynomial combinations of **generators**:

$$\langle \langle 12 \rangle, \langle 23 \rangle \rangle = \left\{ a_1 \langle 12 \rangle + a_2 \langle 23 \rangle \mid a_i \text{ are any spinor polynomials} \right\}.$$

Zariski Nagata Theorem

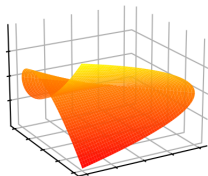
If \mathcal{N} vanishes to order k everywhere where $\langle 12 \rangle = \langle 23 \rangle = 0^*$ then

$$\mathcal{N} \in \langle \langle 12 \rangle, \langle 23 \rangle \rangle^{\langle k \rangle}.$$

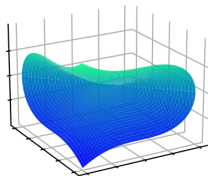
* and $\langle \langle 12 \rangle, \langle 23 \rangle \rangle$ is radical.

Branching of Surfaces Defined by Polynomials

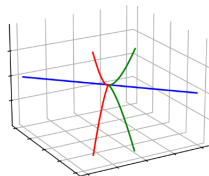
- When we intersect surfaces, we may have **multiple branches**.



$$xy^2 + y^3 - z^2 = 0$$



$$x^3 + y^3 - z^2 = 0,$$



$$xy^2 + y^3 - z^2 = x^3 + y^3 - z^2 = 0.$$

- Our double denominator zero surface has two branches:

$$\langle 12 \rangle = \langle 23 \rangle = 0 \quad \Leftrightarrow \quad \langle 12 \rangle = \langle 23 \rangle = \langle 13 \rangle = 0 \quad \text{or} \quad \lambda_2^\alpha = 0.$$

- We **compute** branchings with primary decomposition techniques.

[De Laurentis, BP '22], see also [Zhang '12].

Ansatz Construction Algorithm, Sketched

- 1 Construct branches of surfaces where **two denominators vanish**.

$$\mathcal{D}_i = \mathcal{D}_j = 0 \quad \longrightarrow \quad \mathcal{V} = \{U_1, U_2, \dots\}.$$

- 2 Sample **near surface** to determine degree of numerator vanishing.

$$U: \mathcal{D}_i \sim \mathcal{D}_j \sim \epsilon \quad \Rightarrow \quad \mathcal{N}_k \sim \epsilon^{\kappa_U}.$$

- 3 Ansatz is basis of intersection of associated **ideals of vanishing polynomials**. Ansatz constructed using **Gröbner basis** techniques.

$$\mathcal{N}_k \in \bigcap_{U \in \mathcal{V}} I(U)^{\langle \kappa_U \rangle}.$$

$$\mathcal{A} = \frac{\mathcal{N}}{\langle 12 \rangle \langle 23 \rangle \langle 34 \rangle [45] [56] [61] s_{345}}.$$

- Probe 108 surfaces where pairs of $\langle ij \rangle$, $[ij]$, s_{ijk} are **small**.

e.g. $[12] \sim [13] \sim [23] \sim \mathcal{O}(\epsilon) \quad \Rightarrow \quad \mathcal{A} \sim \mathcal{O}(\epsilon^2).$

- \mathcal{N} vanishes non-trivially on 28 surfaces. Many **ideal memberships**:

$$\mathcal{N} \in \langle [12], [13], [23] \rangle^2 \cap \langle \langle 12 \rangle, \langle 34 \rangle \rangle \cap \langle \langle 12 \rangle, [16] \rangle \cap (25 \text{ more}).$$

- Imposing that \mathcal{N} is a degree six polynomial gives **one term Ansatz**:

$$\mathcal{N} = c_0 \left(\langle 12 \rangle [21] \langle 45 \rangle [54] \langle 4|2+3|1 \rangle + [16] \langle 6|1+2|3 \rangle \langle 34 \rangle s_{123} \right), \quad c_0 \in \mathbb{Q}.$$

Proof-of-Concept Remainders for $q\bar{q} \rightarrow \gamma\gamma\gamma$

$$\mathcal{A} = \text{Diagram} + \dots$$

(Simulated evaluations using analytics from [Abreu, BP, Pascual, Sotnikov '20]).

- Analyze remainder, reconstruct pentagon function coefficients.
- Reconstruction now requires $317 \mathbb{Q}_p + 566 \mathbb{F}_p$ samples.

Amplitude	$R_{-++}^{(2,0)}$	$R_{-++}^{(2,N_f)}$	$R_{+++}^{(2,0)}$	$R_{+++}^{(2,N_f)}$
Ansatz Dim [Abreu et al '20]	41301	2821	7905	1045
Ansatz Dim [De Laurentis, BP '22]	566	20	18	6

Some Thoughts on Bottlenecks

Ansatz Basis Construction in [2203.04269]

- \mathcal{N} lives in space of polynomials of fixed mass-dimension/little group:

$$\mathcal{N} \in \mathcal{M}_{d, \vec{\phi}}.$$

- Approach used in [2203.04269] is to avoid intersecting ideals:

$$\text{Ansatz} = \text{basis} \left(\bigcap_{U \in \mathcal{V}} \left[\mathcal{M}_{d, \vec{\phi}} \cap I(U)^{\langle \kappa_U \rangle} \right] \right).$$

- Bottlenecks:
 - Many, large vector space intersections.
 - Ansatz picked by RREF. c_i reconstruction needs many finite fields.

$$\mathcal{N} = \sum_i c_i \mathbf{a}_i, \quad (c_i \bmod p_1, c_i \bmod p_2, \dots) \xrightarrow{\text{CRT}} c_i.$$

Improvement path: Can we **analytically intersect ideals** in controlled way?

Regular Sequences

- Ideal intersection is combinatoric when using a **regular sequence**:

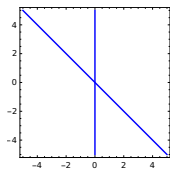
$\{q_1, \dots, q_n\}$ where q_{i+1} is **non-zero divisor** mod $\{q_1, \dots, q_i\}$, $i \in [0, n - 1]$.

- $\{\langle 12 \rangle, \langle 13 \rangle, \langle 14 \rangle\}$ is not a regular sequence because of Schouten.

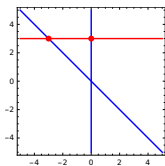
$$\langle 14 \rangle \langle 23 \rangle = \langle 12 \rangle \langle 34 \rangle - \langle 13 \rangle \langle 24 \rangle.$$

- Geometrically, sequential intersections **decrease dimension**.

E.g. $\{x(x+y), y-3\}$ is regular sequence, $\{x(x+y), y-3, x+3\}$ is not:

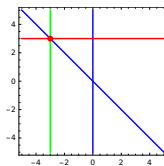


— $x(x+y) = 0$



— $x(x+y) = 0$

— $-3 + y = 0$



— $x(x+y) = 0$

— $-3 + y = 0$

— $x = -3$

Improving Basis Construction

- Regular sequence $\{q_1, \dots, q_n\}$. Monomial ideals intersect as

$$\left\langle \prod_i q_i^{\alpha_i} : \vec{\alpha} \in A \right\rangle \cap \left\langle \prod_i q_i^{\beta_i} : \vec{\beta} \in B \right\rangle = \left\langle \text{lcm} \left(\prod_i q_i^{\alpha_i}, \prod_i q_i^{\beta_i} \right) : \vec{\alpha} \in A, \vec{\beta} \in B \right\rangle. \quad (*)$$

- E.g., $\{\langle 12 \rangle, [12], s_{123}\}$ is regular sequence:

$$\langle \langle 12 \rangle, s_{123}^2 \rangle \cap \langle [12], s_{123} \rangle = \langle \langle 12 \rangle [12], s_{123}^2, \langle 12 \rangle s_{123}, [12] s_{123}^2 \rangle.$$

- **Strategy:**

- 1 Organize intersection into denominator ideals.

$$\bigcap_{U \in \mathcal{V}} I(U)^{\langle \kappa_U \rangle} = \bigcap_{(a,b)} \langle \mathcal{D}_a, \mathcal{D}_b \rangle^{k_{ab}}.$$

- 2 Apply (*) to subsets of intersectands built from regular sequences.

Summary:

- Rational functions in amplitudes have **poorly understood structure**.
- We study them in **singular limits** to characterize that structure. We interpret this behavior in terms of **ideals** from which we build Ansätze.

Looking forward:

- Can we get even better at intersecting these ideals?
(Does there exist a “next-to-regular” sequence?)
- What is the physical origin of the partial fractions structure?

Backup

Lorentz Invariance

- Coefficients are **Lorentz invariant** functions of spinor brackets.

$$\mathcal{C}(\lambda, \tilde{\lambda}) = \mathcal{C}(\langle \rangle, \llbracket \rrbracket).$$

- Relevant ring is Lorentz invariant **subring** of S_n .

$$\mathcal{S}_n = \mathbb{F} \left[\langle 12 \rangle, \dots, \langle (n-1)n \rangle, [12], \dots, [(n-1)n] \right].$$

- Variables are brackets, now have “**Schouten identities**”.

$$\mathcal{J}_{\Lambda_n} = \left\langle \sum_{j=1}^n \langle ij \rangle [jk], \langle ij \rangle \langle kl \rangle - \langle ik \rangle \langle jl \rangle - \langle il \rangle \langle kj \rangle, \langle \rangle \leftrightarrow \llbracket \rrbracket \right\rangle.$$

- Physical **spinor bracket functions** also form a quotient ring.

$$\mathcal{R}_n = \mathcal{S}_n / \mathcal{J}_{\Lambda_n}.$$

Bases of Spinor Space and Polynomial Reduction

- Numerators are \mathbb{Q} -linear combinations of spinor monomials.

$$m_\alpha = \prod_i v_i^{\alpha_i} \quad \text{where} \quad \vec{v} = \{\langle 12 \rangle, \langle 23 \rangle, \dots, [12], [23], \dots\}.$$

- Polynomial reduction writes p in terms of generators g_i .

$$p = \Delta_{\{g_1, \dots, g_k\}}(p) + \sum_{i=1}^k c_i g_i.$$

- Polynomial in ideal if and only if Groebner remainder is 0.

$$\Delta_{\mathcal{G}(J)}(p) = 0 \quad \Leftrightarrow \quad p \in J.$$

- Monomials irreducible by $\mathcal{G}(\mathcal{J}_{\Lambda_n})$ form basis. Related [Zhang '12]

$$\text{basis} = \{m_\alpha \text{ such that } \Delta_{\mathcal{G}(\mathcal{J}_{\Lambda_n})}(m_\alpha) = m_\alpha\}.$$

How To Perform Numerical Investigations?

- Need to find phase-space points $(\lambda^\epsilon, \tilde{\lambda}^\epsilon)$ where \mathcal{D}_i are small.

$$\mathcal{D}_i(\lambda^\epsilon, \tilde{\lambda}^\epsilon) \sim \mathcal{D}_j(\lambda^\epsilon, \tilde{\lambda}^\epsilon) \sim \epsilon.$$

- Conflict with modern techniques: **no small elements** in a finite field.

$$|0|_{\mathbb{F}_p} = 0, \quad \text{and} \quad a \neq 0 \Rightarrow |a|_{\mathbb{F}_p} = 1.$$

- Approaching with complex numbers would be plagued by instabilities.

Enter the **p-adic numbers** – a middle ground between finite fields and \mathbb{C} .

Introduction to the p -adic Numbers

- The p -adic numbers roughly correspond to **Laurent series** in p .

$$x = \sum_{i=\nu}^{\infty} a_i p^i = a_{\nu} p^{\nu} + a_{\nu+1} p^{\nu+1} + \dots, \quad \left(\begin{array}{l} a_i \in [0, p-1], \\ a_{\nu} \neq 0. \end{array} \right).$$

- The p -adic numbers form a **field**. $x, y \in \mathbb{Q}_p \Rightarrow$

$$x + y \in \mathbb{Q}_p, \quad -x \in \mathbb{Q}_p, \quad x \times y \in \mathbb{Q}_p, \quad \frac{1}{x} \in \mathbb{Q}_p \text{ (if } x \neq 0\text{)}.$$

- The p -adic **absolute value** allows for small numbers ($p \sim \epsilon$).

$$|x|_p = p^{-\nu}, \quad \Rightarrow \quad |p|_p < |1|_p.$$

Computing with p -adic Numbers

- For computing purposes* we **truncate to finite order**.

$$x = p^{\nu(x)} \left(\underbrace{\tilde{x}}_{\text{mantissa}} + \mathcal{O}(p^k) \right).$$

*Try [<https://github.com/GDeLaurentis/pyadic>] to investigate yourselves.

- Truncation **reduces to finite field** case for $\nu = 0, k = 1$.
- Arithmetic ($+ - /*$) is essentially performed **modulo p^k** , e.g.

$$x \times y = p^{\nu(x)+\nu(y)} \left(\tilde{x}\tilde{y} + \mathcal{O}(p^k) \right).$$

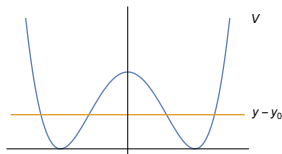
- Mantissa inverse computed with **extended euclidean algorithm**.

p -adic (Integer) Points Near an Irreducible Variety

- Want to find $(\lambda^{(\epsilon)}, \tilde{\lambda}^{(\epsilon)})$ “close” to $U = V(\langle q_1, \dots, q_m \rangle_{R_n})$:

$$q_i(\lambda^{(\epsilon)}, \tilde{\lambda}^{(\epsilon)}) = p c_i + \mathcal{O}(p^k), \quad \sum_{i=1}^n \lambda_{i\alpha}^{(\epsilon)} \tilde{\lambda}_{i\dot{\alpha}}^{(\epsilon)} = 0 + \mathcal{O}(p^k).$$

- First, find finite field $x \in U$ by intersecting with random plane.



- Arbitrarily extend \mathbb{F}_p point $(\lambda, \tilde{\lambda})$ to k digits. Trivially near U .
- To satisfy momentum conservation, perturb by $(p\delta, p\tilde{\delta})$.

$$(\lambda^{(\epsilon)}, \tilde{\lambda}^{(\epsilon)}) = (\lambda + p\delta, \tilde{\lambda} + p\tilde{\delta}).$$

Polynomials that Vanish on a Variety

- Polynomials that vanish on **all points** of U form an ideal

$$I(U) = \left\{ q \in S_n \text{ where } q(x) = 0 \text{ for all } x \in U \right\}.$$

- Consider if \mathcal{N}_i **vanishes to order k_U** on U ,

$$\mathcal{N}_i(x^{(\epsilon)}) = \mathcal{O}(\epsilon^{k_U}), \text{ where } |x - x^{(\epsilon)}| \leq \epsilon \text{ and } x \in U.$$

- It turns out that \mathcal{N}_i **still belongs to an ideal!**

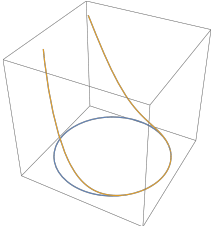
Zariski-Nagata Theorem

Polynomials vanishing to $\mathcal{O}(k_U)$ on U belong to $I(U)^{\langle k_U \rangle}$ – the k_U th “**symbolic power**” of $I(U)$.

- Computed from **primary decomposition** of ideal power $I(U)^{k_U}$.

Examples of Symbolic Powers

- A function **vanishing to fourth order** at a point on the circle:

$$\langle x - 1 \rangle_{\mathbb{F}[x,y]/\langle x^2+y^2-1 \rangle}^{\langle 4 \rangle} \sim \cdot$$


- Often the symbolic power **coincides** with standard power, e.g.

$$\langle \langle 12 \rangle, [12] \rangle_{R_5}^{\langle 2 \rangle} = \langle \langle 12 \rangle, [12] \rangle_{R_5}^2 = \langle \langle 12 \rangle^2, \langle 12 \rangle [12], [12]^2 \rangle_{R_5}.$$

- Symbolic/standard power **may not coincide**. E.g. in $\mathbb{F}[x, y, z]$

$$\langle xy, xz, yz \rangle^{\langle 2 \rangle} = \langle x^2y^2, x^2z^2, y^2z^2, xyz \rangle \neq \langle xy, xz, yz \rangle^2$$

The p -adic Logarithm

- Over the p -adic numbers, one can define **converging power series**.
- The power series for a logarithm converges for $|x|_p < 1$.

$$\log_p(1+x) = \sum_{k=1}^{\infty} \frac{(-1)^{k+1} x^k}{k}$$

- To **map to radius of convergence**, use Fermat's little theorem.

$$w^{p-1} = 1 \pmod{p} \quad \Rightarrow \quad |w^{p-1} - 1|_p < 1$$

- Logarithm relations then p -adically **analytically continue** \log_p .

$$\log_p(w) = \frac{1}{p-1} \log(w^{p-1})$$