

Quantum Computing and Quantum Machine Learning

Michael Spannowsky

IPPP, Durham University

P3H Summer School

Lecture

Michael Spannowsky



The Speed Of Light

The Morning After: Google claims 'quantum supremacy'

And a controversial 'Ghost in the Shell' trailer.



E





Master in Elektrotechnik, Informatik, Robotik, Maschinenwesen o. ä. (w/m/d)

German Aerospace Center (DLR) · Oberpfaffenhofen, Bavaria, Germany (On-site)

4 company alumni

Professor Cyber Security im Online Fernstudium (m/w/d)

IU International University of Applied Sciences · Germany (Remote)

Actively recruiting



Actively recruiting

Master Thesis: Design of digitally enhanced power management circuits for Future Quantum Computers

Forschungszentrum Jülich · Jülich, North Rhine-Westphalia, Germany (On-site)

1 company alum



Expertin für Quantenkommunikation (w/m/d)

Deutsche Bahn · Frankfurt, Hesse, Germany (On-site)

Actively recruiting



Kristen Philipkoski

Published 10 years ago: September 2, 2011 at 7:02 am - Filed to: COMPUTING \sim





Share 😝 🗲 🕁

A OMPUTING

hris Ferrie and whurley

uantum Computers Will Be Incredibly Useful For

Computers don't exist in a vacuum. They serve to solve problems, and the type of problems they can solve are influenced by their hardware. Graphics processors are specialized for rendering images; artificial intelligence processors for AI; and quantum computers designed for... what? While the power of quantum computing is impressive, it does not mean that existing ...





"Nature is quantum [...] so if you want to simulate it, you need a quantum computer" – Richard Feynman (1982)

Easily said ... so how do we do that?

Beginning of a scientific journey that accelerated in recent years tremendously....

Pre-digital age computing:

- express values in combinations and positions of beads
- manipulate beads mechanically
- convert position and combination of beads back into value



Digital age computing:



Quantum Age computing



We then measure one specific outcome. Have to repeat measurement to statistically evaluate how likely each outcome is (by calculating and measuring several times). Since we work only with probabilities, we measure only probabilities

How can these quantum principles help to improve computations?

classical system is in one state out of 16			quantum (superposition) can be in all states at same time				
			0011	0000	0001	0010	0011
			0111	0100	0101	0110	0111
	1001		1011	1000	1001	1010	1011
			1111	1100	1101	1110	1111

- Configuration space here 16=2⁴ states.
- Computations can be performed simultaneously on the whole configuration space. -> can be much faster than classically
- A measurement of the quantum system after the computations are performed results in the observation of one of these configurations, with a probability that corresponds to the computational processes

6

Quantum Gate



quantum gate and multi slit experiment are conceptually identical

It's a secret computation...

While operating one cannot see how the gate works. Only at the end one can measure the outcome (box is closed during operations)



Galton Board as analogy for Quantum Computer



The quantum mechanical principles on which the algorithms have to rely to have a chance for a quantum advantage are



P3H Summer School

Lecture

What are the potential advantages of Quantum Devices?

- Quantum Simulation: Simulating genuine quantum systems, e.g. molecules Medical applications, Chemistry applications, HEP etc
- Quantum Cryptography and Encryption and decryption, e.g. emails, RSA Security:
- Quantum Information Science:
 Transformation, storage and transmission of information, e.g. databases, teleportation, networks etc
- Quantum Sensing and Quantum Metrology: Imaging, Quantum Navigation/Timing
- Quantum Machine Learning Learning and optimisation based and Optimisation:
 On quantum algorithms
- Speed and efficiency: Simple classical tasks performed faster or with less electricity consumption etc

P3H Summer School Lecture

Private and Public Sector is placing big bets on Quantum Computing

Quantum Computing Use Cases





Gartner



Significant financial investment expected across many sectors

In US, already now higher financial investment from private than public sector



All national and international labs have QC programmes (Fermilab, BNL, LBNL, CERN, ...)

P3H Summer School

Lecture

Michael Spannowsky

October 2024

Basic motivation for Quantum Computing

"Can we take the quantum mechanical properties of microscopic objects and scale them up to larger quantum systems while harnessing their quantum prowess?"



For some specialised task quantum supremacy has been shown

Disclaimer:

nobody today thinks that quantum computers will universally replace classical computers Technical challenges of a quantum computer

• Many quantum paradigms require system to be perfectly isolated (shielded from outside) to maintain coherence – for as long as the algorithm takes



The road to Quantum Advantage



Lecture

14

P3H Summer School



Qubit technologies overview. From: Forbes, <u>Quantum Computer</u> <u>Battle Royale: Upstart Ions Versus Old Guard Superconductors</u>



Quantum Information & Physics



P3H Summer School

Lecture



Figure 2. Quantum computing prototypes announced on vendor roadmaps

Source: Arthur D. Little, Olivier Ezratty

Lecture

16

Analog vs Digital Quantum Computing

Analog and digital quantum computing are two different approaches to quantum computing, each with its own advantages and disadvantages.

Analog Quantum Computing (AQC):

- Based on the principle of quantum evolution of a quantum system, e.g. quantum annealing
- The system uses its intrinsic quantum dynamics, following the Schroedinger Equation
- Ground state represents the solution to the problem at hand
- Not always universal, but often well-suited for optimisation problems

Example: D-Wave Systems. The D-Wave quantum annealer uses a network of qubits that can collectively tunnel through the solution space to find the global minimum of a given function.

Digital Quantum Computing (DQC):

- Digital quantum computing, also known as gate-based quantum computing
- Uses quantum logic gates to perform operations on qubits
- Considered to be more versatile than analog computing.
- However, might require higher level of control over the quantum system, which can be challenging

Example: IBM's and Google's quantum computers use the gate-based model of quantum computing.

Popular Quantum Computing paradigms

Quantum computing has long and distinguished history but is only now becoming practicable.

Туре	Discrete Gate (DG)	Continuous Variable (CV)	Quantum Annealer (QA)
Property	Universal (any quantum algorithm can be expressed)	Universal - GBS non-Universal	Not universal — certain quantum systems
Advantage	most algorithms and tech support	uncountable Hilbert (configuration) space	continuous time quantum process
How?	IBM - Qiskit ~ 100 Qubits	Xanadu	DWave - LEAP ~7000 Qubits
What?			
	input 10) -H - Z - H - C - C - C - C - C - C - C - C - C	Interferometer U_1 S U_2 D Φ	CA finds wide region failed tunnelling state

Lecture





Quantum Mechanics Basics





Lecture

Phenomenological observations:

• Randomness of measurement outcomes:

Repeated measurements of the same physical quantity (observable) A in the same physical conditions (state) produce different results.

Post-measurement state: Let ψ be the physical state of the considered quantum system. If we perform a measurement process on the system to measure the observable A and the obtained outcome is a ∈ R then the state of the system, after the measurement, is ψa.

• Incompatible observables: There are pairs of observables that cannot be simultaneously measured by an experiment.

Mathematical description of QM for QC

• Define and work in **Hilbert space**

- Physical states are elements of Hilbert space
- States are manipulated through linear operators $A(\alpha\psi + \beta\varphi) = \alpha A\psi + \beta A\varphi$
- Physical quantities are expressed through selfadjoint operators, that have real eigenvalues $\lambda \langle \psi | \psi \rangle = \langle \psi | A \psi \rangle = \langle A \psi | \psi \rangle = \lambda^* \langle \psi | \psi \rangle$

- Superposition principle
- Distance and similarity measure
- composition and transformation of objects



- The eigenvalues of unitary operator are complex numbers with unit modulus, also called phases: for φ eigenvector of unitary operator U with eigenvalue μ , $\langle \varphi | \varphi \rangle = \langle U \varphi | U \varphi \rangle = \mu^* \mu \langle \varphi | \varphi \rangle = |\mu|^2 \langle \varphi | \varphi \rangle$, then $|\mu|^2 = 1$, so $\mu = e^{i\varphi}$ with $\varphi \in \mathbb{R}$
- Spectral Theorem: If A ∈ B(H) is normal, that is, AA⁺ = A⁺A, if and only if there exists an orthonormal basis of H made by eigenvectors of A.

 - ----> Functional calculus for bounded (and unbounded) operators

• Composite quantum systems:

Let HA and HB be Hilbert spaces and $\psi \in$ HA, $\phi \in$ HB.

The tensor product of ψ and φ is defined by: $\psi \otimes \varphi(x, y) := \langle \psi | x \rangle_A \langle \varphi | y \rangle_B$ $x \in H_A, y \in H_B$ The tensor product of Hilbert spaces $H_A \otimes H_B$ consists of all such tensor products equipped with the inner product: $\langle \psi \otimes \varphi | \psi' \otimes \varphi' \rangle := \langle \psi | \psi' \rangle_A \cdot \langle \varphi | \varphi' \rangle_B$ $\psi, \psi' \in H_A \varphi, \varphi' \in H_B$

The tensor product of operators $A \in B(HA)$ and $B \in B(HB)$ is: $(A \otimes B)(\psi \otimes \varphi) := A\psi \otimes B\varphi$

If dim HA = n and dim HB = m then dim(HA \otimes HB) = n·m.

For example:

• for
$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle$$
 and $|\phi\rangle = \gamma |0\rangle + \delta |1\rangle$ we get
 $|\psi\rangle \otimes |\phi\rangle = \alpha \gamma |0\rangle \otimes |0\rangle + \alpha \delta |0\rangle \otimes |1\rangle + \beta \gamma |1\rangle \otimes |0\rangle + \beta \delta |1\rangle \otimes |1\rangle$

• for operators $A|\psi\rangle=a|\psi\rangle$ and $B|\phi\rangle=b|\phi\rangle$ we have:

 $(A \otimes B)(|\psi\rangle \otimes |\phi\rangle) = (A|\psi\rangle) \otimes (B|\phi\rangle) = a|\psi\rangle \otimes b|\phi\rangle = ab(|\psi\rangle \otimes |\phi\rangle)$

•
$$A = \begin{pmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & \cdots & a_{nn} \end{pmatrix} \quad B = \begin{pmatrix} b_{11} & \cdots & b_{1m} \\ \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \vdots \\ b_{m1} & \cdots & b_{mm} \end{pmatrix} \rightarrow A \otimes B = \begin{pmatrix} a_{11}B & \cdots & a_{1n}B \\ \vdots & \vdots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1}B & \cdots & a_{nn}B \end{pmatrix}$$

P3H Summer School Lecture 23 Michael Spannowsky October 2024

• Some states cannot be written as a direct tensor product, e.g.

$$|\psi\rangle_{AB} = |\psi\rangle_A \otimes |\psi\rangle_B \quad \neq \quad \frac{1}{\sqrt{2}}|0\rangle_A|0\rangle_B + \frac{1}{\sqrt{2}}|1\rangle_A|1\rangle_B$$
 Bell states

So, using the language of states and their tensor products is not be sufficient.

$$|\psi_x\rangle \longrightarrow \rho = \sum_x p_x |\psi_x\rangle \langle \psi_x$$

Density matrix:

Consider a quantum system with state space C^d . A density matrix, commonly denoted as ρ , is a linear operator $\rho \in L(C^d, C^d)$ such that: 1. $\rho \ge 0$, and 2. $tr(\rho) = 1$.

The density matrix is a representation of a system's statistical state

• Concepts pure vs mixed states, coherent vs incoherent superposition

• Entanglement

The density operator $\rho \in S(HA \otimes HB)$ is said to be separable if it can be written as a statistical mixture of product states: $\sum (A) = (B)$

$$\rho = \sum_{i} \lambda_i \rho_i^{(A)} \otimes \rho_i^{(B)}$$

where $\lambda_i \ge 0$ and $\sum_i \lambda_i = 1$. Otherwise it is said to be **entangled**.

Example:
$$|\Psi
angle=rac{1}{\sqrt{2}}(|00
angle+|11
angle)$$
 Bell state

Assume the Bell state could be written into product of two states

 $|\Psi\rangle = |a\rangle \otimes |b\rangle = (\alpha|0\rangle + \beta|1\rangle) \otimes (\gamma|0\rangle + \delta|1\rangle) = \alpha\gamma|00\rangle + \alpha\delta|01\rangle + \beta\gamma|10\rangle + \beta\delta|11\rangle$



Let's reconsider the phenomenological evidence list:

Randomness of measurement outcomes:

The possible experimental values of the observable A are the element of its spectrum $\sigma(A)$.

Given a pure state $|\psi\rangle \in H$, the probability of measuring the value $a \in \sigma(A)$ is:

$$\mathbb{P}_{\psi}(a) = \langle \psi | P_a \psi \rangle$$

where $\{P_a\}_{a\in\sigma}(A)$ is the spectral measure or projection value measure (PVM) of A. If we consider repeated measurements of the observable A in the same physical conditions represented by the state $|\psi\rangle$, the expectation value of A on the state $|\psi\rangle$ is the mean of the outcomes:

$$\langle A \rangle_{\psi} := \sum_{a \in \sigma(A)} a \mathbb{P}_{\psi}(a) = \langle \psi | A \psi \rangle$$

Post-measurement state: Let $|\psi\rangle \in H$ be the state of the considered quantum system. If we perform a measurement process of A with outcome $a \in \sigma(A)$, then the state of the system, after the measurement, is:

$$|\psi_a\rangle = \frac{P_a|\psi\rangle}{\sqrt{\langle\psi|P_a\psi\rangle}}$$

In the presented mathematical formulation of quantum mechanics the measurement process of the observable A is completely described by the PVM $\{P_a\}_{a\in\sigma}(A)$ which determines the probability distribution of the outcomes and the post-measurement state.

Compatible and incompatible observables: A and B are compatible when they commute:

$$[A,B] := AB - BA = 0$$

in this case: $P_a^A P_b^B = P_b^B P_a^A \forall a \in \sigma(A)$ and $\forall b \in \sigma(B)$, so the following probability is well-defined:

$$\mathbb{P}_{\psi}(A = a \land B = b) = \langle \psi | P_a^A P_b^B \psi \rangle = \langle \psi | P_b^B P_a^A \psi \rangle$$

P3H Summer School Lecture 27 Michael Spannowsky October 2024

Rotation about the Bloch Sphere and state parametrisation





Extending this to a system of N qubits forms a 2^N -dimensional Hilbert Space

Quantum dynamics

The time evolution of an isolated quantum system is mathematically described by a one-parameter group of unitary operators $\{U_t\}_{t\in R}$ defined by:

$$U_t := \sum_{\lambda \in \sigma(\mathcal{H})} e^{-i\frac{t}{\hbar}\lambda} P_h \equiv e^{-i\frac{t}{\hbar}\mathcal{H}}$$

where \hbar is the reduced Planck constant, H is the Hamiltonian operator which represents the observable total energy of the considered system and $\{P_h\}_{h\in\sigma}(H)$ is the spectral measure of H.

If the state at time t = 0 is $|\psi_0\rangle \in H$ then the state at time t > 0 is:

$$|\psi_t\rangle = U_t |\psi_0\rangle = e^{-i\frac{t}{\hbar}\mathcal{H}} |\psi_0\rangle$$

Taking the time derivative obtains the Schroedinger equation:

$$i\hbar \frac{d}{dt} |\psi_t\rangle = \mathcal{H} |\psi_t\rangle$$

In case of a time-dependent Hamiltonian, H must be replaced by a one-parameter family of self-adjoint operators $\{H(t)\}_{t\in R}$ and the Schroedinger equation assumes the form:

$$i\hbar \frac{d}{dt} |\psi_t\rangle = \mathcal{H}(t) |\psi_t\rangle$$

P3H Summer School

Lecture

Michael Spannowsky

Hamiltonian simulation

Recall Schroedinger Equation and time evolution equation

$$|\psi(t)\rangle = \hat{U}(t) |\psi(0)\rangle \rightarrow \frac{\partial \hat{U}(t)}{\partial t} = -iH\hat{U}(t) \rightarrow \hat{U}(t) = e^{-iHt}$$

If $H = H^{\dagger}$ (Hermitian) then $\hat{U}(t)$ is unitary $e^{-iHt}e^{iHt} = 1$ and $e^{-iHt} = (e^{iHt})^{\dagger}$

$$e^{-iHt} = \sum_{j=0}^{\infty} \frac{(-iHt)^j}{j!}$$



[2111.00627] [2301.00560]

Pauli Operators

Tensor products of Pauli operators {I, X, Y, Z} form a basis for the vector space of 2ⁿ x 2ⁿ complex matrices, which are used to represent quantum states and operators in a system of n qubits. Thus, our hermitian matrix H can be decomposed into Pauli operators

Trotterisation

For commuting self-adjoint operators [S,T] = 0 we find

$$e^{S+T}\xi = e^S e^T \xi, \quad \xi \in D(S) \cap D(T)$$

However, if S and T do not commute this doesn't hold. Surprisingly the Trotter Product Formula comes to the rescue:

$$\operatorname{s-lim}_{n \to \infty} \left(e^{-i\frac{t}{n}S} e^{-i\frac{t}{n}T} \right)^n = e^{-it(S+T)}$$

Consequently

$$e^{A+B} \approx \left(e^{A/N}e^{B/N}\right)^N \quad \approx \left[\left(I + \frac{A}{N}\right)\left(I + \frac{B}{N}\right)\right]^N = \left[I + \frac{A}{N} + \frac{B}{N} + \frac{AB}{N^2}\right]^N \quad \approx \left[I + \frac{A+B}{N}\right]^N$$

- If N too large causes numerical instabilities, but must be sufficiently large
- Trotterization error, important error for quantum algorithms (Hamiltonian simulation, time evolution etc)

P3H Summer School Lecture 31 Michael Spannowsky October 2024

Trotterization tells us the error we make when writing ${\cal H}$ as a sum of ${\cal H}_i$

$$H = H_1 + H_2 + H_3 + \cdots + H_N$$
 for $U(H, t) = e^{-iHt/\hbar}$

Thus we implement an approximated time evolution where the H_i are compositions of Pauli matrices

$$U(H,t,n) \;=\; \prod_{j\,=\,1}^n \prod_k e^{-i H_k t/n} \qquad H \;=\; \sum_k H_k,$$

Each piece remains unitary, and H_i is hermitian

$$H = H_1 + H_2 + \cdots + H_N$$

$$= U(H_1) = U(H_2) = \cdots = U(H_N) =$$

Task is to convert each piece into gate operations

P3H	Summer	School	
-----	--------	--------	--



Trotterization error needs to be assessed, e.g. by reducing time steps

From QM to QFT

- Extend QM to systems with variable particle numbers (quantum many-body problems)
- Promote classical fields to operator-valued functions, acting on states in Fock space
- Second Quantisation Programme

• Time Evolution Operator: $|\Psi(t)\rangle = \hat{U}(t,t_0)|\Psi(t_0)\rangle, \quad \hat{U}(t,t_0) = \mathcal{T}\exp\left(-\frac{i}{\hbar}\int_{t_0}^t \hat{H}(t')dt'\right)$

perturbative approach

e.g. scatting in weak coupling regime

- Interaction picture (split Hamiltonian into free and interaction H)
- define Dyson series: perturbative expansion, suitable if coupling is small
- ➡ Wick contractions, Feynman diagrams
- ➡ S-Matrix, LSZ theorem

non-perturbative approach

e.g. Real-time time evolution suitable for large couplings

One usually works in the Schroedinger picture (states are time-dependent and operators are time-independent (unless they are explicitly time-dependent))

 $\mathcal{F} = \bigoplus^{\sim} \mathcal{H}^{\otimes n}$

➡ Latticisation, Kogut-Susskind programme

2

Why Hamiltonian simulation? \rightarrow The infamous sign problem

- Sign problem profound challenge for simulation of field theories
- Can arise in presence of chemical potential, topological terms, multiparticle dynamics, ...
- Example chemical potential $\mu \bar{\psi} \gamma^0 \psi$

$$Z = \int \mathscr{D}\bar{\psi} \mathscr{D}\psi \mathscr{D}A \ e^{-S[\bar{\psi},\psi,A]} \quad \text{(partition function)}$$

$$S = \int_0^{1/T} d\tau \int d^3x \left[\bar{\psi} (\gamma^\mu D_\mu + m) \psi + \frac{1}{4} F^a_{\mu\nu} F^{a\mu\nu} + \mu \bar{\psi} \gamma^0 \psi \right]$$

and integration over fermion fields and Wick rotation (imaginary time)

$$Z = \int \mathscr{D}Ae^{-S_{\text{gauge}}[A]} \cdot \det(\gamma^{\mu}D_{\mu} + m + \mu\gamma^{4}) \quad \longrightarrow \text{ For } \mu \neq 0 \text{ complex phases don't cancel}$$



The infamous sign problem

- Importance sampling Interpretation of $e^{-S_{gauge}} \det(M)$ as probability weight
- Highly oscillatory integrands

 $\langle O \rangle = \frac{\int \mathscr{D}Ae^{-S_{\text{gauge}}} O |\det[M(A)]| e^{i\phi(A)}}{\int \mathscr{D}Ae^{-S_{\text{gauge}}} |\det[M(A)]| e^{i\phi(A)}}$

near cancellation of pos and neg contribs



 $\int dx \exp(-x^2 + i\lambda x) \to \int dx \exp(-x^2) \cos(\lambda x)$



[de Forcrand '10]

P3H Summer School

Lecture

Michael Spannowsky

HEP application focused quantum simulations

->

 Real-time evolution on quantum computer can avoid sign problem

 \bullet Continuous field theories $\phi(x)$ describe particle phenomenology

$$|\psi(0)\rangle = e^{-\frac{i}{\hbar}\hat{H}t} = |\psi(t)\rangle$$

infinite dimensional `matrices' $\left| \left\langle X(T) \right| U(T, -T) \right| pp(-T) \right\rangle \right|^2$

 $H_{\scriptscriptstyle B}^{\scriptscriptstyle {
m Field}}$

 Needs discretisation irrespective of classical or quantum computation



Lecture

Steps to Hamiltonian Simulation on the lattice (Kogut-Susskind)

We consider a non-Abelian gauge theory with fermionic matter fields. The gauge group is general and denoted as SU(N)

The Lagrangian density is given by
$$\mathcal{L} = -\frac{1}{4}F^a_{\mu\nu}F^{\mu\nu a} + \bar{\psi}\left(i\gamma^\mu D_\mu - m\right)\psi$$

38

where:

 ψ : Fermion fieldm: Fermion mass γ^{μ} : Gamma matrices satisfying Clifford algebrastructure constant
of gauge group $F^{a}_{\mu\nu}$: Field strength tensor $F^{a}_{\mu\nu} = \partial_{\mu}A^{a}_{\nu} - \partial_{\nu}A^{a}_{\mu} + gf^{abc}A^{b}_{\mu}A^{c}_{\nu}$ D_{ψ} : Covariant derivative $D_{\mu}\psi = (\partial_{\mu} - igT^{a}A^{a}_{\mu})\psi$
generators of

gauge group

Deriving the Hamiltonian density

$$\begin{array}{ll} \mbox{Gauge fields } A^{a}_{\mu} : & \pi^{a\mu} = \frac{\partial \mathcal{L}}{\partial(\partial_{0}A^{a}_{\mu})} = -F^{0\mu a} & \mbox{electric field components} \\ & \mbox{spatial components} & \pi^{ai} = -F^{0ia} = E^{ia} \\ & \mbox{temporal component } (\mu = 0) & \pi^{a0} = 0 \end{array}$$

Fermion fields
$$\psi, \bar{\psi}$$
: $\pi_{\psi} = \frac{\partial \mathcal{L}}{\partial(\partial_0 \psi)} = i\psi^{\dagger}$ $\pi_{\bar{\psi}} = \frac{\partial \mathcal{L}}{\partial(\partial_0 \bar{\psi})} = 0$

 $\bar{\psi}$ does not have a canonical momentum associated with it -> primary constraint thus, it is not an independent dynamical variable Hamlitonian needs to be constructed carefully to take constraint into account

Legendre transformation:

Hamiltonian density
$$\mathcal{H} = \pi^{a\mu} \partial_0 A^a_\mu + \pi_\psi \partial_0 \psi - \mathcal{L}$$

P3H Summer School

Lecture

replace the ∂_0 component in terms of conjugate momenta (Hamilton approach)

Gauge field:
$$\pi^{ai} = -F^{0ia} = -\left(\partial^0 A^{ia} - \partial^i A^{0a} + gf^{abc}A^{0b}A^{ic}\right) \longrightarrow \partial^0 A^{ia} = -\pi^{ai} - D^i A^{0a}$$

with $D^i A^{0a} = \partial^i A^{0a} - af^{abc}A^{ib}A^{0c}$

Fermion field: $\partial_0 \psi = -i \left(\alpha \cdot D - \beta m \right) \psi - g A_0^a T^a \psi$ with Dirac matrices α and β



Discretisising Space: Introducing the lattice

Spatial lattice: Lattice sites defined by integer coordinates $\mathbf{n} = (n_x, n_y, n_z)$

Lattice spacing defined by the value a (fixed distance between neighbouring sites)

→ serves as UV cutoff

The time remains continuous to retain the Hamiltonian formulation

Link variables: When discretising a gauge theory onto a lattice, we replace the continuous spacetime with a discrete set of points (sites). The gauge fields $A_{\mu}(x)$, which live on continuous space-time, need to be represented in a way that preserves gauge invariance on the lattice

-> variables U are introduced such that they maintain gauge invariance

defined as $U_i(\mathbf{n}) = e^{iagA_i^a(\mathbf{n})T^a}$

The link variable $U_i(\mathbf{n})$ represents the **parallel transporter** (also known as the **Wilson line**) along the link from site \mathbf{n} to site $\mathbf{n} + \hat{i}$. It encodes the phase factor acquired by a particle moving through the gauge field along that link.

Fermion Fields: $\psi(\mathbf{n})$ placed at the lattice sites, representing matter fields at each space point

Define plaquette operator $U_{\mu\nu}(\mathbf{n}) = U_{\mu}(\mathbf{n})U_{\nu}(\mathbf{n}+\hat{\mu})U_{\mu}^{\dagger}(\mathbf{n}+\hat{\nu})U_{\nu}^{\dagger}(\mathbf{n})$

The plaquette operator is a measure of the curvature (field strength) of the gauge field over the area of the plaquette. In the limit of small lattice spacing a, $U_{\mu\nu}(\mathbf{n})$ approximates the exponential of the field strength tensor integrated over the plaquette area.

For small
$$a$$
 we have $U_{\mu}(\mathbf{n}) = e^{iagA_{\mu}^{a}(\mathbf{n})T^{a}} \simeq 1 + iagA_{\mu}^{a}(\mathbf{n})T^{a} - \frac{a^{2}g^{2}}{2}(A_{\mu}^{a}(\mathbf{n})T^{a})^{2} + \mathcal{O}(a^{3})$

Calculating the plaquette for small a

with $A_{\nu}^{b}(\mathbf{n}+\hat{\mu}) = A_{\nu}^{b}(\mathbf{n}) + a\partial_{\mu}A_{\nu}^{b}(\mathbf{n}) + \mathcal{O}(a^{2})$

$$U_{\mu\nu}(\mathbf{n}) = \left(1 + iagA^a_{\mu}(\mathbf{n})T^a\right) \left(1 + iagA^b_{\nu}(\mathbf{n} + \hat{\mu})T^b\right)$$
$$\times \left(1 - iagA^c_{\mu}(\mathbf{n} + \hat{\nu})T^c\right) \left(1 - iagA^d_{\nu}(\mathbf{n})T^d\right) + \mathcal{O}(a^3)$$
$$= 1 + ia^2gF^a_{\mu\nu}(\mathbf{n})T^a + \mathcal{O}(a^3)$$

With the plaquette we can express the latticised magnetic energy term ${\cal H}_{\cal B}$ as the trace over the plaquete:

$$H_B = \frac{1}{g^2} \sum_{\mathbf{n}, i < j} \left(N_c - \operatorname{Re} \operatorname{Tr} \left[U_{ij}(\mathbf{n}) \right] \right) \approx \frac{a^4}{2} \sum_{\mathbf{n}, i < j} F_{ij}^a(\mathbf{n}) F_{ij}^a(\mathbf{n})$$

gauge invariant quantity

Electric energy term $H_E = \frac{g^2}{2} \sum_{\mathbf{n},i,a} [E_i^a(\mathbf{n})]^2$

Fermion energy term $H_F = H_K + H_M$

Lecture

P3H Summer School

Kinetic term:
$$H_{K} = \frac{1}{2a} \sum_{\mathbf{n},i} \left[\psi^{\dagger}(\mathbf{n}) \alpha^{i} U_{i}(\mathbf{n}) \psi(\mathbf{n}+\hat{i}) - \psi^{\dagger}(\mathbf{n}+\hat{i}) \alpha^{i} U_{i}^{\dagger}(\mathbf{n}) \psi(\mathbf{n}) \right]$$

Mass term:
$$H_{M} = m \sum_{\mathbf{n}} \psi^{\dagger}(\mathbf{n}) \beta \psi(\mathbf{n})$$

Ensure gauge invariance of H on the lattice group element at site \mathbf{n} Local gauge transformations at each site \mathbf{n} $\psi(\mathbf{n}) \to G(\mathbf{n})\psi(\mathbf{n})$ and Link $U_i(\mathbf{n}) \to G(\mathbf{n})U_i(\mathbf{n})G^{\dagger}(\mathbf{n}+\hat{i})$

-> Electric energy term
$$[E_i^a(\mathbf{n})]^2$$
 is gauge invariant

- \rightarrow Trace of plaquette variable $Tr[U_{ij}(\mathbf{n})]$ is gauge invariant
- The combination $\psi^{\dagger}(\mathbf{n})U_{i}(\mathbf{n})\psi(\mathbf{n}+\hat{i})$ is gauge invariant

The discrete version of Guass's law operator is $G^{a}(\mathbf{n}) = \sum_{i} \left[E_{i}^{a}(\mathbf{n}) - E_{i}^{a}(\mathbf{n} - \hat{i}) \right] + g\psi^{\dagger}(\mathbf{n})T^{a}\psi(\mathbf{n})$ Physical Hilbert space states must satisfy $G^{a}(\mathbf{n})|\Psi\rangle = 0$

43

Michael Spannowsky

October 2024

The fermion doubling problem

Consider $H_F = H_K + H_M$ for the free theory (i.e. no gauge interactions) $U_i(\mathbf{n}) = 1$

Fourier transforming $\psi(\mathbf{n}) = \int_{-\frac{\pi}{a}}^{\frac{\pi}{a}} \frac{d^d p}{(2\pi)^d} e^{i\mathbf{p}\cdot\mathbf{n}a} \tilde{\psi}(\mathbf{p})$, inserting into H_K and massaging the

equation gives

$$H_{K} = \int \frac{d^{d}p}{(2\pi)^{d}} \tilde{\psi}^{\dagger}(\mathbf{p}) \left[\frac{2i\alpha^{i}}{a} \sin(ap_{i}) \right] \tilde{\psi}(\mathbf{p}) = \int \frac{d^{d}p}{(2\pi)^{d}} \tilde{\psi}^{\dagger}(\mathbf{p}) \ (\alpha \cdot \mathbf{K}) \ \tilde{\psi}(\mathbf{p})$$

where $\mathbf{K} = \frac{2}{a}(\sin(ap_1), \sin(ap_2), \sin(ap_3))$

-> the eigenvalues are given by $E(\mathbf{p}) = \pm |\mathbf{K}| = \pm \frac{2}{a} \sqrt{\sum_{i=1}^{d} \sin^2(ap_i)}$

the energy vanishes when $\sin(ap_i) = 0$, which occurs at $p_i = 0$ and $p_i = \frac{\pi}{a}$

Thus, there are 2^d points in the Brillouin zone where $E(\mathbf{p}) = 0$. In four dims (d=4) there are 16 such points, indicating 16 fermion species (including doublers)

Introducing staggered fermions

To mitigate the fermion doubling problem, we introduce staggered fermions, which involve:

- Replacing the multi-component Dirac spinor $\psi(\mathbf{n})$ with a single-component fermion field $\chi(\mathbf{n})$
- Redistributing the spinor components across neighbouring lattice sites using staggered phases $\eta_i(\mathbf{n}) = (-1)^{n_1+n_2+\ldots+n_{i-1}}$

The staggered fermion Hamiltonian (for $U_i(\mathbf{n}) = 1$ for simplicity) is

$$H_F = \frac{1}{2a} \sum_{\mathbf{n},i} \eta_i(\mathbf{n}) \left[\chi^{\dagger}(\mathbf{n}) \chi(\mathbf{n}+\hat{i}) - \chi^{\dagger}(\mathbf{n}+\hat{i}) \chi(\mathbf{n}) \right] + m \sum_{\mathbf{n}} \chi^{\dagger}(\mathbf{n}) \chi(\mathbf{n})$$

After Fourier transform, one has

$$H_F = \int \frac{d^d k}{(2\pi)^d} \,\tilde{\chi}^{\dagger}(\mathbf{k}) \left[\frac{2i}{a} \sum_i \sin(ak_i) \,\tilde{\eta}_i \right] \tilde{\chi}(\mathbf{k}) \,+\, m \int \frac{d^d k}{(2\pi)^d} \,\tilde{\chi}^{\dagger}(\mathbf{k}) \tilde{\chi}(\mathbf{k})$$

By distributing the spinor components across different lattice sites and introducing the staggered phases, the number of fermion species is reduced

- in d dimensions, the number of species reduces from 2^d to $2^{d/2}$
- In four dimensions, from 16 to 4

While this doesn't completely eliminate the fermion doubling problem, it significantly reduces the number of unphysical doublers.

	P3H Summer School	Lecture	45	Michael Spannowsky	October 2024
--	-------------------	---------	----	--------------------	--------------

Assembling the Lattice Hamiltonian

 $H = H_F + H_R + H_F$ The total Hamiltonian is $H_E = \frac{g^2}{2} \sum_{\mathbf{n},i,a} \left[E_i^a(\mathbf{n}) \right]^2$ Electric Energy Term $H_B = \frac{1}{g^2} \sum_{\mathbf{n}, i < j} \left[N_c - \operatorname{Re} \operatorname{Tr} \left(U_{ij}(\mathbf{n}) \right) \right]$ Magnetic Energy Term $H_F = \frac{1}{2a} \sum_{\mathbf{n},i} \eta_i(\mathbf{n}) \left[\chi^{\dagger}(\mathbf{n}) U_i(\mathbf{n}) \chi(\mathbf{n}+\hat{i}) - \chi^{\dagger}(\mathbf{n}+\hat{i}) U_i^{\dagger}(\mathbf{n}) \chi(\mathbf{n}) \right]$ Fermion Energy Term (Staggered Fermions) $+m\sum_{\mathbf{n}}\chi^{\dagger}(\mathbf{n})\chi(\mathbf{n})$ H_{B}^{Field} Numerical methods for Hamiltonian simulation: H_E^{Field} H_{v}^{int} • Tensor Network Approches Quantum Simulations H^{Matter} $H_h^{\rm int}$

46

Concrete example U(1) in 1-dimension

The U(1) gauge field is continuous $U(n) = e^{i\theta(n)}$ and has to be truncated

- → Assume 2 sites
- --- Truncate the Link Variables U(n) e.g. assume $\theta(n) = 0, \frac{2\pi}{3}, \frac{4\pi}{3}$

Truncate fermions – assume single fermion mode per site, i.e. each site either 0 or 1 fermions

3 states for electric field on each site, thus 3 links => $3 \times 3 = 9$ $|-1, -1\rangle, |-1, 0\rangle, |-1, +1\rangle, |0, -1\rangle, |0, 0\rangle, |0, +1\rangle, |+1, -1\rangle, |+1, 0\rangle, |+1, +1\rangle$ each site 0 or 1 fermions x 2 sites => fermion Hilber space $2^2 = 4$ $|0,0\rangle, |0,1\rangle, |1,0\rangle, |1,1\rangle$ Hamiltonian 36×36

Thus, 36 basis states are labeled as $|E_0, E_1; n_0, n_1\rangle$ Calculate elements $H_{ij} = \langle E'_0, E'_1; n'_0, n'_1 | H | E_0, E_1; n_0, n_1 \rangle$ with $H = H_E + H_K + H_M$ (H_B is absent in 1d) diagonal element $H_{ii} = \frac{1}{2}(E_0^2 + E_1^2)$ off-diagonal element $H_{ij} = -\frac{i}{2}\eta(0)U(E_0) \ \delta_{E'_0,E_0-1} \ \delta_{n'_0,n_0-1} \ \delta_{n'_1,n_1+1}$

P3H Summer School

Lecture

47

Michael Spannowsky

October 2024

Hamiltonian simulation – what resources do we need?

• Discretisation of field $\phi_n(x_k)$

• On quantum devices algorithms require exp less resources $\ln_2[\dim H]$

Quantum computing not optional for Hamiltonian simulation Effective Field Theories can ameliorate problem

P3H Summer School

Lecture

48

Michael Spannowsky

Quantum Circuits

Lecture

Need transition form classical to quantum:

Single-Qubit Quantum Gates

Illustrative to write single-qubit operation as matrices

X-Gate: Quantum equivalent to classical NOT gate

 $\begin{array}{c} |0\rangle \mapsto |1\rangle \\ |1\rangle \mapsto |0\rangle \end{array}$

 $\rightarrow \text{ Flips |0> to |1> and vice versa (hopping)}$ Represented by matrix $\mathbf{X} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ concretely $\mathbf{X}|0\rangle = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} = |1\rangle$

It is unitary $\mathbf{X}\mathbf{X}^{\dagger} = \mathbf{X}\mathbf{X}^{-1} = \mathbb{1}$

Z-Gate: Represented by matrix
$$\begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

Action $|0\rangle \mapsto |0\rangle$ $|1\rangle \mapsto -|1\rangle$

Note, the X, Y and Z gates are represented by the Pauli matrices

$$\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad \sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \quad \sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \quad [\sigma_i, \sigma_j] = 2i\epsilon_{ijk}\sigma_k$$

P3H Summer School Lecture 52 Michael Spannowsky October 2024

gate: Matrix representation
$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

Hadamard

$$h: \quad |0\rangle \mapsto \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \iff |+\rangle := \frac{|0\rangle + |1\rangle}{\sqrt{2}}$$
$$|1\rangle \mapsto \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \iff |-\rangle := \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$

Phase gate: Matrix representation
$$P_{\phi} := \begin{pmatrix} 1 & 0 \\ 0 & e^{i\phi} \end{pmatrix}$$

With special phase values

$$S := P_{\pi/2}$$
 $T := P_{\pi/4}$ $R := P_{-\pi/4}$

Summary of fixed 1-qubit gates:

Gate	Circuit representation	Matrix representation	Dirac representation
X	- <u>X</u> -	$ \left(\begin{array}{c} 0 & 1\\ 1 & 0 \end{array}\right) $	$ 1\rangle\langle 0 + 0\rangle\langle 1 $
Y	- <u>Y</u> -	$ \left(\begin{array}{cc} 0 & -i \\ i & 0 \end{array}\right) $	$i 1\rangle\langle 0 -i 0\rangle\langle 1 $
Ζ	- <u>Z</u> -	$ \left(\begin{array}{cc} 1 & 0\\ 0 & -1 \end{array}\right) $	$ 1\rangle\langle 0 - 0\rangle\langle 1 $
Н	-H	$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1\\ 1 & -1 \end{pmatrix}$	$\frac{1}{\sqrt{2}}(0\rangle + 1\rangle)\langle 0 + \frac{1}{\sqrt{2}}(0\rangle - 1\rangle)\langle 1 $
S	_ <u>S</u>	$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}$	$\frac{1}{\sqrt{2}} 0\rangle\langle 0 + \frac{1}{\sqrt{2}}i 1\rangle\langle 1 $
Т	-T	$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 \\ 0 & e^{(-i\pi/4)} \end{pmatrix}$	$\frac{1}{\sqrt{2}} 0\rangle\langle 0 + \frac{1}{\sqrt{2}}e^{(-i\pi/4)} 1\rangle\langle 1 $

Quantum gate can be parametrised

Pauli rotations:

$$R_{x}(\theta) = e^{-i\frac{\theta}{2}\sigma_{x}} = \begin{pmatrix} \cos\left(\frac{\theta}{2}\right) & -i\sin\left(\frac{\theta}{2}\right) \\ -i\sin\left(\frac{\theta}{2}\right) & \cos\left(\frac{\theta}{2}\right) \end{pmatrix} = \cos\frac{\theta}{2}I - i\sin\frac{\theta}{2}X$$
$$R_{y}(\theta) = e^{-i\frac{\theta}{2}\sigma_{y}} = \begin{pmatrix} \cos\left(\frac{\theta}{2}\right) - \sin\left(\frac{\theta}{2}\right) \\ \sin\left(\frac{\theta}{2}\right) & \cos\left(\frac{\theta}{2}\right) \end{pmatrix} = \cos\frac{\theta}{2}I - i\sin\frac{\theta}{2}Y$$
$$R_{z}(\theta) = e^{-i\frac{\theta}{2}\sigma_{z}} = \begin{pmatrix} e^{-i\frac{\theta}{2}} & 0 \\ 0 & e^{i\frac{\theta}{2}} \end{pmatrix} = \cos\frac{\theta}{2}I - i\sin\frac{\theta}{2}Z$$

generalised form via $R(\theta_1, \theta_2, \theta_3) = R_z(\theta_1)R_y(\theta_2)R_z(\theta_3)$

$$R(\theta_1, \theta_2, \theta_3) = \begin{pmatrix} e^{i(-\frac{\theta_1}{2} - \frac{\theta_3}{2})} \cos(\frac{\theta_2}{2}) & -e^{i(-\frac{\theta_1}{2} + \frac{\theta_3}{2})} \sin(\frac{\theta_2}{2}) \\ e^{i(\frac{\theta_1}{2} - \frac{\theta_3}{2})} \sin(\frac{\theta_2}{2}) & e^{i(\frac{\theta_1}{2} + \frac{\theta_3}{2})} \cos(\frac{\theta_2}{2}) \end{pmatrix}$$

P3H Summer School Lecture 55

Measurement process

Measurement process of a generic (normalised) qubit state $|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle$

represented by projection onto eigenstates $P_0 = |0\rangle\langle 0|$ and $P_1 = |1\rangle\langle 1|$

Prob of measurement outcome 0 is then $p(0) = tr(P_0|\psi\rangle\langle\psi|) = \langle\psi|P_0|\psi\rangle = |\alpha_0|^2$

and
$$p(1) = |\alpha_1|^2$$

After measurement qubit is in state

$$|\psi\rangle \leftarrow \frac{P_0|\psi\rangle}{\sqrt{\langle\psi|P_0|\psi\rangle}} = |0\rangle$$

The observable corresponding to a computational basis measurement is Pauli-Z observable

 $\sigma_z = |0\rangle\langle 0| - |1\rangle\langle 1|$ (we know eigenvalues +1 for |0> and -1 for |1>)

The expectation value $\langle \sigma_z \rangle$ in a value in [-1, 1]. Its error can be estimated as sampling from a Bernoulli distribution.

Wald interval gives

→ For $\epsilon = 0.1$ and conf level 99% one needs 167 samples For $\epsilon = 0.01$ and conf level 99% one needs 17,000 samples

→ Overall might need a large number of shots on quantum computer This needs to be taken into account when comparing quantum and classical computers in terms of speedups and quantum advantage

The Bloch Sphere

Since
$$|\psi
angle=lpha\,|0
angle+eta\,|1
angle$$
 with $|lpha|^2+|eta|^2=1$ one can find angles such that

$$lpha = oldsymbol{e}^{i\gamma}\cosrac{ heta}{2} \qquad eta = oldsymbol{e}^{i\delta}\sinrac{ heta}{2}$$

Thus, with $\phi = \delta - \gamma$ single qubit can be parametrised as

$$|\psi\rangle = e^{(i\gamma)} \left(\cos\frac{\theta}{2}|0\rangle + e^{(i\phi)}\sin\frac{\theta}{2}|1\rangle\right)$$

where a global imaginary phase has no measurable effect and can be omitted.

 $(\sin\theta\cos\phi,\sin\theta\sin\phi,\cos\phi)$

2-qubit states

Are built by tensor products, each qubit can be in state |0> or in state |1> So, for two qubits we have four possibilities:

```
\left|0\right\rangle \otimes \left|0\right\rangle, \left|0\right\rangle \otimes \left|1\right\rangle, \left|1\right\rangle \otimes \left|0\right\rangle, \left|1\right\rangle \otimes \left|1\right\rangle
```

that we denote

```
\left|0\right\rangle \left|0\right\rangle ,\left|0\right\rangle \left|1\right\rangle ,\left|1\right\rangle \left|0\right\rangle ,\left|1\right\rangle \left|1\right\rangle
```

or

$$\ket{00}, \ket{01}, \ket{10}, \ket{11}$$

We can have superposition as a generic state

Lecture

$$|\psi
angle = lpha_{00} |00
angle + lpha_{01} |01
angle + lpha_{10} |10
angle + lpha_{11} |11
angle$$

with complex coefficients such that $\sum_{x,y=0}^{1} |lpha_{xy}|^2 = 1$

59

2-qubit states

Furthermore, we can express the state as a vector

 $\begin{pmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{10} \\ \alpha_{11} \end{pmatrix}$

For which we find the inner products

$$\langle 00|00\rangle = \langle 01|01\rangle = \langle 10|10\rangle = \langle 11|11\rangle = 1$$

$$\langle 00|01\rangle = \langle 00|10\rangle = \langle 00|11\rangle = \cdots = \langle 11|00\rangle = 0$$

A 2-qubit quantum gate is a unitary matrix U of size 4 x 4

2-qubit gates

CNOT gate: unitary matrix representation $\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$

> In words: if the first qubit is |0> nothing changes. If it is |1> we flip the second bit (and first stays the same)

Action: $|00\rangle \rightarrow |00\rangle$ $|01\rangle \rightarrow |01\rangle$ $|10\rangle \rightarrow |11\rangle$ $|11\rangle \rightarrow |10\rangle$ As a gate: $x, y \in \{0, 1\}$ \rightarrow $|x\rangle$ -- $|y\rangle$ -- $|y\rangle$ $|y \oplus x\rangle$

- A set of gates that can approximate any quantum operation
 -> Universal quantum computer
 - e.g. Rotation gates $R_x(\theta), R_y(\theta), R_z(\theta)$ + phase shift gate $P(\varphi)$ + CNOT

The CNOT gate is an extremely important gate

- It realises conditional probabilities
- It creates entanglement

Bell state (fully entangled)

• It can copy classical information, because

|00
angle
ightarrow |00
angle

```
|10
angle 
ightarrow |11
angle
```

Constructs other control gates

P3H Summer School Lecture 62 Michael Spannowsky October 2024

N-qubit states

When we have n qubits, each of them can be in state 10> or 11>

Thus for n qubit states we have 2ⁿ possibilities:

$$|00\ldots0\rangle,|00\ldots1\rangle,\ldots,|11\ldots1\rangle$$

or simply

$$\left|0\right\rangle,\left|1\right\rangle,\ldots,\left|2^{n}-1\right\rangle$$

A generic state of the system will be

$$|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle + \ldots + \alpha_{2^n-1} |2^n-1\rangle$$

With complex coefficients, such that

$$\sum_{i=0}^{2^{n}-1} |\alpha_{i}|^{2} = 1$$

Suppose we have the N qubit state

$$|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle + \ldots + \alpha_{2^n-1} |2^n-1\rangle$$

If we measure all its qubits, we obtain:

- 0 with probability $|lpha_0|^2$ and the new state will be $|0\dots00
 angle$
- 1 with probability $|\alpha_1|^2$ and the new state will be $|0...01\rangle$
- ...
- $2^n 1$ with probability $|\alpha_{2^n-1}|^2$ and the new state is $|1 \dots 11\rangle$

Completely analogous to 1 and 2 qubit situation but now with 2^n possibilities

Example: Turning a Hamiltonian term into a gate

Recall
$$H = H_1 + H_2 + \cdots + H_N$$

$$= U(H_1) = U(H_2) = \cdots = U(H_N) = U(H_N)$$

Assume, universal gate operations on device are $\{H, R_Z, CX\}$

Example 1 Assume
$$H_1 = Z \longrightarrow U = e^{-iZt} \longrightarrow R_Z(2t)$$

 $R_Z(\theta) = e^{-i\frac{\theta}{2}Z}$

Assume $H_2 = X \longrightarrow$ Since $HXH = Z \Rightarrow X = HZH$ Example 2

65

 $\longrightarrow U = He^{-iZt}H$ (proof via CBH Formula)

$$\rightarrow$$
 H $R_Z(2t)$ H

Example 3 $H = Z \otimes Z$ note $e^{-Z \otimes Zt} \neq e^{-iZt} \otimes e^{-iZt}$

with
$$(Z \otimes Z)^2 = \mathbb{I}$$
 one finds $e^{i(Z \otimes Z)t} = \cos(t)\mathbb{I} - i\sin(t)Z \otimes Z$

for the action on states we find

$$e^{i(Z \otimes Z)t} |00\rangle = (\cos(t)\mathbb{I} - i\sin(t)Z \otimes Z) |00\rangle = (\cos(t) - i\sin(t)) |00\rangle$$
$$e^{i(Z \otimes Z)t} |11\rangle = (\cos(t)\mathbb{I} - i\sin(t)Z \otimes Z) |11\rangle = (\cos(t) - i\sin(t)) |11\rangle$$
$$e^{i(Z \otimes Z)t} |01\rangle = \cos(t) |01\rangle - i\sin(t)Z |0\rangle \otimes Z |1\rangle = (\cos(t) + i\sin(t)) |01\rangle$$

which can be written in matrix form as

$$e^{i(Z\otimes Z)t} = \begin{bmatrix} e^{-it} & 0 & 0 & 0\\ 0 & e^{it} & 0 & 0\\ 0 & 0 & e^{it} & 0\\ 0 & 0 & 0 & e^{-it} \end{bmatrix}_{|1\rangle}^{|0\rangle} \text{ if } \# \text{ of } 1 \text{ is even one gets -} (\text{parity of state})$$

$$(\text{parity of state})$$

$$(\text{$$

The Ising Model

- The Ising model is a fundamental mathematical model in statistical mechanics used to understand phase transitions and critical phenomena, particularly in ferromagnetic materials.
- Originally proposed by Wilhelm Lenz 1920, extensively studied by his student Ernst Ising 1925
- The model considers a lattice of spins in one of two states: up (+1) or down (-1). These spins represent magnetic dipole moments of atomic spins in material and interact with their nearest neighbours. -> Ferromagnetism, critical phenomena and phase transitions
- Exact solutions in 1-D and 2-D (Onsager's solution 1944) no exact solution in 3-D (NP-hard)

$$\begin{array}{ll} \text{interaction strength J} & \text{magnetisation} \\ \\ \text{Classical Ising Model:} & H = -\sum_{\langle i,j \rangle} J_{i,j} \, \sigma_i \sigma_j - \mu \sum_i h_i \, \sigma_i \\ & \text{sum over neighbouring spins} \\ \\ \text{Quantum Mechanical Ising Model:} & H = -J \sum_{\langle i,j \rangle} \sigma_i^z \sigma_j^z - h \sum_i \sigma_i^z - \Gamma \sum_i \sigma_i^x \\ \\ \text{Real-time evolution } U(t) = e^{-iHt} \text{ with } H = -J \sum_{i=1}^{N-1} \sigma_i^z \sigma_{i+1}^z - \Gamma \sum_{i=1}^N \sigma_i^x \\ \end{array}$$

Lecture

67

Michael Spannowsky

Now that the model has been constructed and Trotterised, we are ready to implement the real-time evolution on a quantum device:

3-site Ising model - one trotter step:

	•		.	
P3H Summer School	Lecture	68	Michael Spannowsky	October 2024

Examples of HEP usecases for Hamiltonian simulation

- Real-time evolution in QFTs: scattering processes, quark-gluon plasma formation, out-of-equilibrium dynamics
- Sign-problem in finite density QCD: QCD phase diagram at finite baryon density or nuclear matter in neutron stars
- Simulating Early Universe Physics: Phase transitions, reheating
- Neutrino Oscillations in Dense Media: Neutrino osciallations in supenovae, neutron stars or early universe environments
- Gauge theories in higher dimensions:
- Topological QFTs:

P3H Summer School

Lecture

69

Simons-like terms in SM Michael Spannowsky October 2024

Chern-Simons theory and Chern-

Classically expensive. Extra dim models

- Quantum computing is a new computational paradigm with a high potential for computational improvements in many science areas
- Hamiltonian simulation is an active research area ideally suited to be executed on quantum devices
- It might be key to avoiding the so-called `sign problem' and to obtain a quantum advantage in computations for fundamental physics
- Hands-on session:

https://github.com/simon-j-williams/QCIsingModel_KIT/tree/main